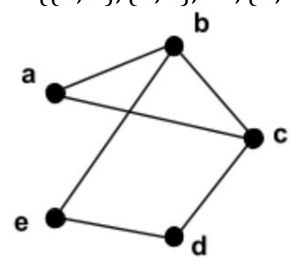
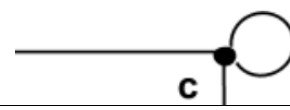
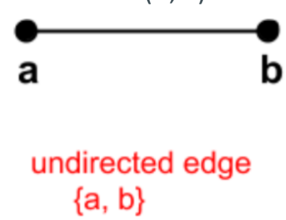
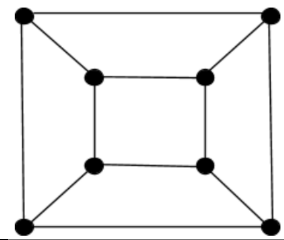
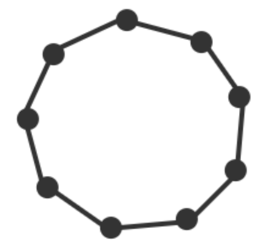
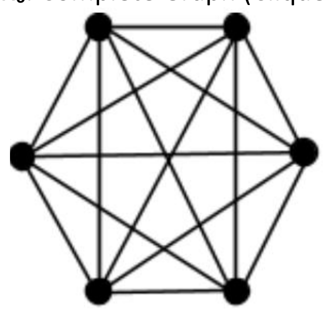
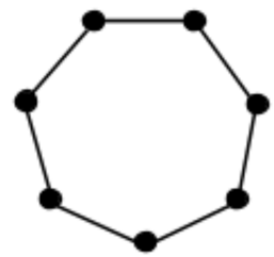
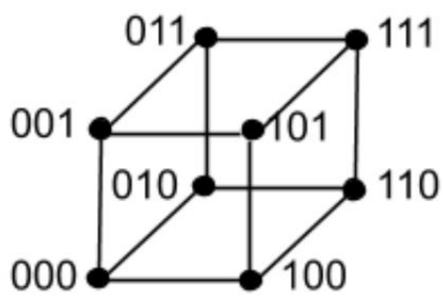
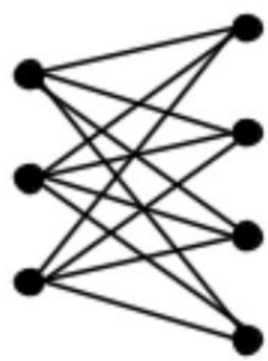
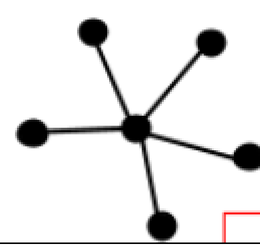
Harold's Undirected Graphs and Trees
Cheat Sheet

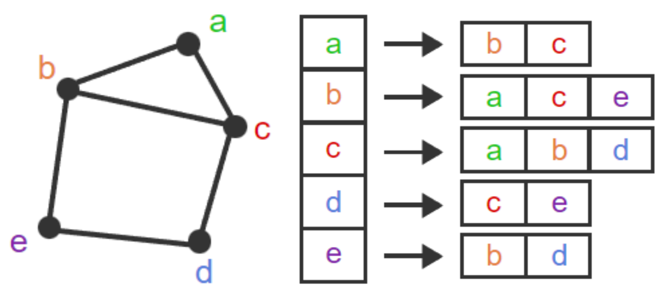
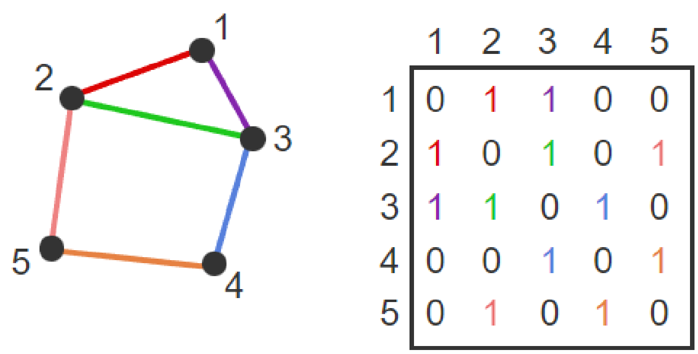
22 October 2022

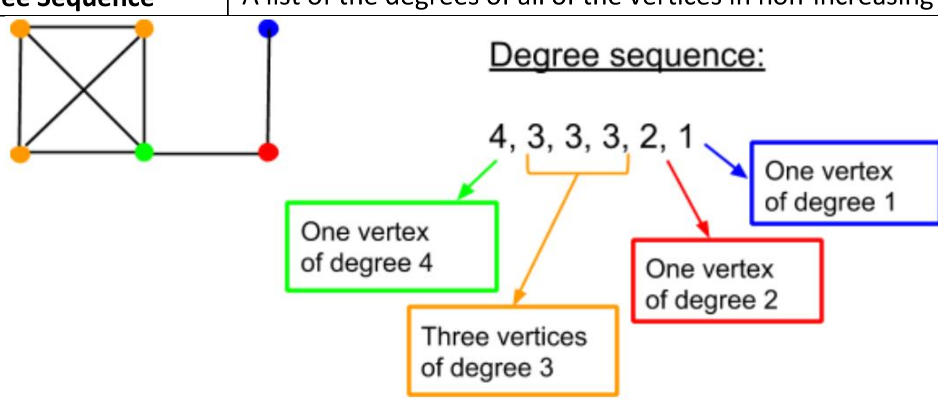
Definitions

Term	Definition	Example
Vertices (Nodes)	An individual element of V is called a vertex. A graph is **finite** if the vertex set is finite.	Set $V = \{a, b, c, d, e\}$ ① or ●
Edges (Arcs)	An edge (u, v) ∈ E, is pictured as an arrow going from one vertex to another.	Set E ⊆ V x V $E = \{\{a, b\}, \{a, c\}, \ldots, \{d, e\}\}$ 
Self-Loop (Loop)	An edge that connects a vertex to itself.	
Undirected Graph	A graph whose edges are unordered pairs of vertices.	G = (V, E) 
Simple Graph	A graph with no parallel edges or self-loops.	|Cycle| ≥ 3
Adjacent	There is an edge between two vertices.	Two vertices are connected.
Endpoints	Vertices b and e are the **endpoints** of edge {b, e}	The two vertices of an edge.
Incident	The edge {b, e} is **incident** to vertices b and e.	The edge of two vertices.
Neighbor	A vertex c is a **neighbor** of vertex b if and only if {b, c} is an edge.	Has an edge to it.
Degree	The **degree** of a vertex is the number of neighbors it has.	$deg(v)$
Total Degree	The sum of the degrees of all of the vertices.	$\sum_{v \in V} deg(v) = 2 \cdot \|E\|$
Regular Graph	All the vertices have the same degree.	$deg(a) = deg(b) = deg(c) \ldots$

d-Regular Graph	All the vertices have degree d.	3-Regular Graph: 
Subgraph	A graph $H = (V_H, E_H)$ is a ***subgraph*** of a graph $G = (V_G, E_G)$ if $V_H \subseteq V_G$ and $E_H \subseteq E_G$. Any graph G is a subgraph of itself.	2-Regular Graph: 
Common Graphs	**K_6**: Complete Graph (Clique)  Has an edge between every pair of vertices.	**C_7**: Cycle 
	Q_3: 3-Dimentional Hypercube  Has 2^n vertices.	**$K_{3,4}$**:  No edges between vertices in the same set.
	P_5: A path 	**S_5**: Star 

Graph Representation

Term	Description
Adjacency List	Each vertex has a list of all its neighbors. 
Matrix	A ‘1’ means an edge is present. Is symmetrical about the diagonal ($M_{i,j} = M_{j,i}$). 
Isomorphic	There is a one-to-one correspondence between each of the edges of two graphs (bijection).
"Efficient" Algorithm	An algorithm that runs in worst-case polynomial time.
Theorem: Vertex degree preserved under isomorphism	Consider two graphs, G and G'. Let f be an isomorphism from G to G'. For each vertex v in G, the degree of vertex v in G is equal to the degree of vertex f(v) in G'. If one graph has a vertex of degree 1 and the other graph does not, then not isomorphic.
Degree Sequence	A list of the degrees of all of the vertices in non-increasing order.



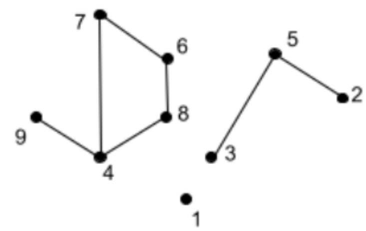
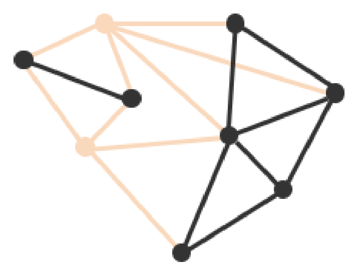
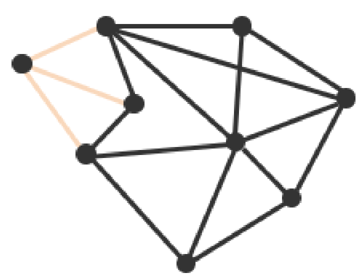
Theorem: Degree sequence preserved under isomorphism	Degree sequence is preserved under isomorphism. If two graphs are isomorphic, they have the same degree sequence.
Graph Theory	Concerned with properties of graphs that are preserved under isomorphism. <u>Preserved</u>: • Number of vertices (|V|) • Number of edges (|E|) • Degree sequence (degrees listed high to low) • Total degree (2·|E|) <u>Not Preserved</u>: • The lowest numbered vertex has degree 3 • Every even numbered vertex has odd degree

Graph Types

Term	Description	Example
Walk	A sequence of alternating vertices and edges that starts and ends with a vertex.	$\langle v_0, (v_0, v_1), v_1, (v_1, v_2) v_2, \ldots, v_l \rangle$ $\langle v_0, v_1, v_2, \ldots, v_l \rangle$ $\lvert \langle v_0 \rangle \rvert = 0$
Open Walk	A walk in which the first and last vertices are not the same.	$\langle a, \ldots, z \rangle$
Closed Walk	A walk in which the first and last vertices are the same.	$\langle a, \ldots, a \rangle$
Length	*l*, the number of edges in the walk, path, or cycle.	$l = \lvert E \rvert$ $l = \lvert V \rvert - 1$ if sequence
Trail	An <u>open</u> walk in which no <u>edge</u> occurs more than once.	$\langle a, b, c, d, c, b, a \rangle$
Circuit	A <u>closed</u> walk in which no <u>edge</u> occurs more than once.	$\langle a, b, a, c, a \rangle$
Path	A trail in which no <u>vertex</u> occurs more than once.	$\langle a, b, c, d \rangle$
Cycle	A circuit of length at least 1 in which no <u>vertex</u> occurs more than once, except the first and last vertices which are the same.	$\langle a, b, c, a \rangle$

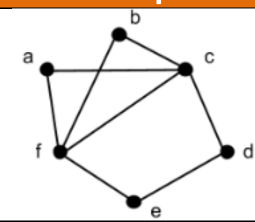
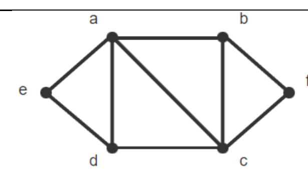
Walk	No repeated …	
	Edge	**Vertex or Edge**
Open	Trail	Path
Closed	Circuit	Cycle

Connectivity

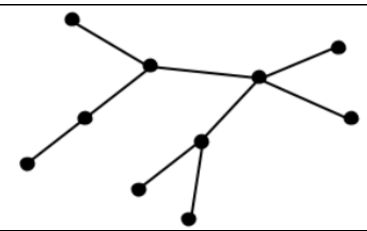
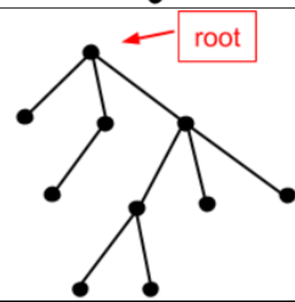
Term	Description	Example
Connected	If there is a path from vertex v to vertex w, then there is also a path from w to v. The two vertices, v and w, are said to be <u>connected</u>.	
Disconnected	A graph is said to be connected if every pair of vertices in the graph is connected, and is <u>disconnected</u> otherwise.	
Connected Component	A maximal set of vertices that is connected.	See graph above for examples.
Isolated Vertex	A vertex that is not connected with any other vertex is called an <u>isolated vertex</u> and is therefore a connected component with only one vertex.	•
k-Vertex-Connected	The graph contains at least k + 1 <u>vertices</u> and remains connected after any k - 1 vertices are **removed** from the graph. (mesh network)	2-vertex-connected: 
Vertex Connectivity	The largest k such that the graph is k-vertex-connected.	κ(G) $\kappa(K_n) = n - 1$
k-Edge-Connected	The graph remains connected after any k - 1 <u>edges</u> are removed from the graph.	3-edge-conncted: 
Edge Connectivity	The largest k such that the graph is k-edge-connected.	λ(G) $\lambda(K_n) = n - 1$
Theorem: Upper bound for vertex and edge connectivity	Let G be an undirected graph. Denote the minimum degree of any vertex in G by *δ(G)*. Then $\kappa(G) \le \delta(G)$ and $\lambda(G) \le \delta(G)$.	The minimum degree of any vertex is at least an upper bound for both the edge and vertex connectivity of a graph.
Complete Graph	There is no set of vertices whose removal disconnects the graph.	Full mesh network.

Euler Circuits and Trails

<table>
<tr><th>Term</th><th>Description</th><th>Example</th></tr>
<tr><td>Euler Circuit</td><td>An undirected graph circuit that contains <u>every</u> edge and <u>every</u> vertex.

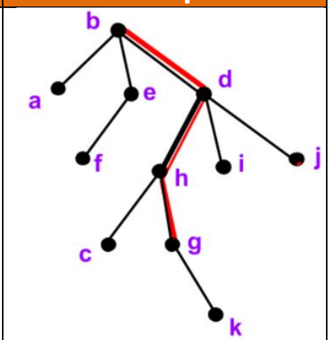
Every vertex reached.
Every edge occurs exactly once.</td><td>
</td></tr>
<tr><td>Theorem: Required conditions for an Euler circuit in a graph</td><td>If an undirected graph G has an Euler circuit, then G is 1) connected and 2) every vertex in G has an even degree.</td><td>$deg(v) = 2k$
where $k \in \mathbb{Z}^+$</td></tr>
<tr><td>Theorem: Sufficient conditions for an Euler circuit in a graph</td><td colspan="2">If an undirected graph G is connected and every vertex in G has an even degree, then G has an Euler circuit.</td></tr>
<tr><td>Theorem: Characterization of graphs that have an Euler circuit</td><td colspan="2">An undirected graph G has an Euler circuit if and only if G is connected and every vertex in G has even degree.</td></tr>
<tr><td>Procedure</td><td colspan="2">Find circuit C in G.
Repeat until C is an Euler circuit:
    Create new graph G' :
        Remove edges in C from G
        Remove isolated vertices
    Find vertex w in G' and C (select any)
    Find circuit C' in G' starting at w
    Combine C and C'
        Follow edges in C to w
        Follow edges in C' back to w
        Follow remaining edges in C
    Rename new circuit to be C</td></tr>
<tr><td>Euler Trail</td><td>An undirected graph open trail that includes <u>each</u> edge exactly once.</td><td>
</td></tr>
<tr><td>Theorem: Characterizations of graphs that have an Euler trail</td><td>An undirected graph G has an Euler trail if and only if G is 1) connected and 2) has exactly <u>two</u> vertices with <u>odd</u> degree.</td><td>Euler trail begins and ends with vertices of odd degree.</td></tr>
</table>

Tree Terms

Term	Description	Example
Tree	An undirected graph that is connected and has no cycles.	Computer file system
Free Tree	There is no particular organization of the vertices and edges	
Rooted Tree	The vertex at the top is designated as the **root** of the tree.	
Level	The **level** of a vertex is its distance (number of edges in the shortest path between the two vertices) from the root.	The root is the only level 0 vertex.
Height	The **height** of a tree is the highest level of any vertex.	Most hops to bottom.
Parent	The first vertex after v encountered along the path from v to the root. (One vertex above v.)	The parent of vertex g is h.
Child	The vertex below the parent.	Vertices c and g are the children of vertex h.
Ancestor	All vertices up in path.	The ancestors of vertex g are h, d, and b.
Descendant	All vertices down in path.	The descendants of vertex h are c, g, and k.
Leaf	Rooted: A vertex which has no children. Free: A vertex of degree 1.	The leaves are a, f, c, k, i, and j. $deg(v) = 1$
Sibling	Vertices with the same parent.	Vertices h, i, and j are siblings of parent d.
Subtree	A tree consisting of new root v and all v's descendants.	The subtree rooted at h includes h, c, g, and k and the edges between them.
Game Tree	Shows all possible playing strategies of both players in a game. Games can be deterministic (tic-tac-toe) or chance (dice).	v_i = game configuration
Variable Length Codes	The number of bits for each character can vary.	'a' = 1, 'e' = 01, etc.

Prefix Code	The code for one character cannot be a prefix of the code for another character.	Leaf nodes guarantee the prefix property.
ASCII	8-Bit characters (256 max.)	UTF-8
Unicode	16-Bit characters (64K max.)	UTF-16
Internal Vertex	<u>Free</u>: The vertex has degree at least two.	$deg(v) \geq 2$
Forest	A graph that has no cycles and that is not necessarily connected. $\|E\| = \|V\| - \|C\|$ (connected components)	

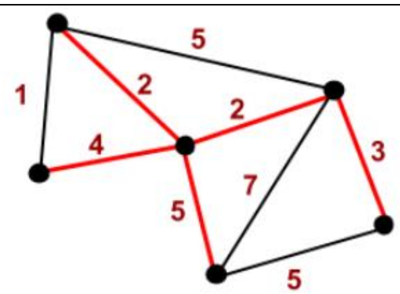
Tree Theorems

Term	Description	Example
Theorem: Unique <u>paths</u> in trees	Let T be a tree and let u and v be two vertices in T. There is <u>exactly one path</u> between u and v. There is a unique path between every pair of vertices in a tree.	
Theorem: Number of <u>edges</u> in a tree	Let T be a tree with n vertices and m edges, then m = n - 1.	$m = n - 1$
Theorem: Number of <u>leaves</u> in a tree	Any free tree with at least two vertices has at least two leaves.	Lower bound 
Theorem: Prim's Algorithm	Prim's algorithm finds a minimum spanning tree of the input weighted graph.	See Spanning Trees below

Tree Traversals

Term	Description	Example
Traversal	Systematically visiting each vertex.	Hit a node.
Pre-Order Traversal	A vertex is visited before its descendants.	First hit (left side) of tree vertex
In-Order Traversal	A vertex is visited after its first descendant.	2^{nd} hit of tree vertex
Post-Order Traversal	A vertex is visited after its descendants.	Last hit (right side) of tree vertex

Spanning Trees

<table>
<tr><th>Term</th><th>Description</th><th>Example</th></tr>
<tr><td>Spanning Tree</td><td>For a connected graph G. a subgraph of G which contains all the vertices in G and is a tree.</td><td>Fewest edges possible to visit all vertices</td></tr>
<tr><td>Depth-First Search (DFS)</td><td>Favors going deep into the graph.
Produces trees with longer paths.</td><td>Explorer ventures far away from home</td></tr>
<tr><td>Breadth-First Search (BFS)</td><td>Explores the graph by distance from the initial vertex, starting with its neighbors and expanding the tree to neighbors of neighbors.
Produces trees with shorter paths.</td><td>Explorer ventures close to home</td></tr>
<tr><td>Weighted Graph</td><td>A graph G = (V ,E), along with a function w: E → ℝ.</td><td>The function w assigns a real number to every edge.</td></tr>
<tr><td>Weight w(G)</td><td>
</td><td>w(G) is the sum of the weights of the edges in G.</td></tr>
<tr><td>Minimum Spanning Tree (MST)</td><td>A spanning tree T of G whose weight is no larger than any other spanning tree of G.</td><td>Goal: Min. weight</td></tr>
<tr><td rowspan="2">Prim's Algorithm</td><td>A classic algorithm for finding minimum spanning trees developed by mathematician Robert Prim in 1957.</td><td>Always choose min. edge in queue.</td></tr>
<tr><td colspan="2"><pre>Input: An undirected, connected, weighted graph G.
Output: T, a minimum spanning tree for G.

T := Ø.
Pick any vertex in G and add it to T.

For j = 1 to n-1
      Let C be the set of edges with one endpoint
inside T and one endpoint outside T.
      Let e be a minimum weight edge in C.
      Add e to T.
      Add the endpoint of e not already in T to T.
End-for</pre></td></tr>
</table>

Sources:

- SNHU MAT 230 - Discrete Mathematics, zyBooks.
- See also “Harold’s Directed Graphs Cheat Sheet”.

